

# Lecture 11 - June 10

## Lexical Analysis

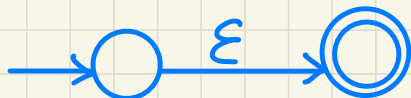
***Formulation: Ext. Subset Construction  
Minimizing DFA***

# Regular Expression to epsilon-NFA

## Base Cases

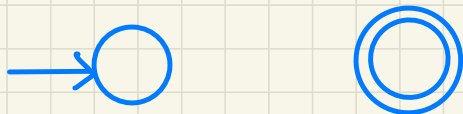
$\epsilon$

$L(\epsilon) = \{\epsilon\}$



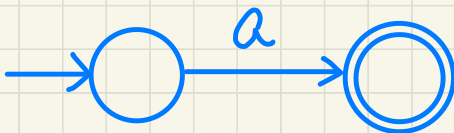
$\emptyset$

$L(\emptyset) = \emptyset$



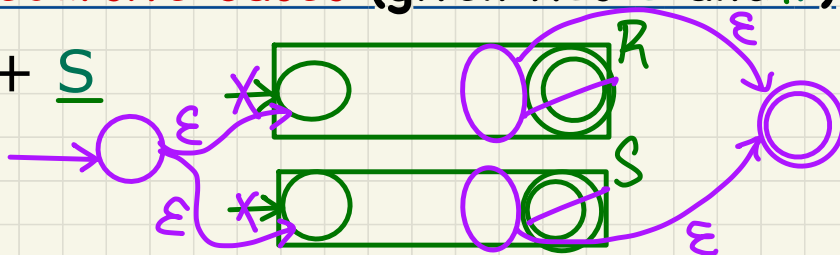
$a$

$(a \in \Sigma)$



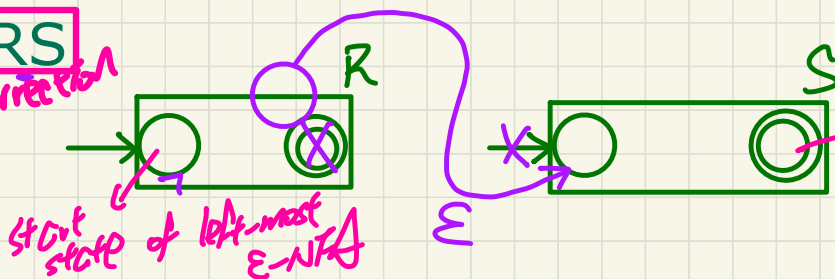
## Recursive Cases (given REs $R$ and $S$ )

$R + S$



$RS$

Correction

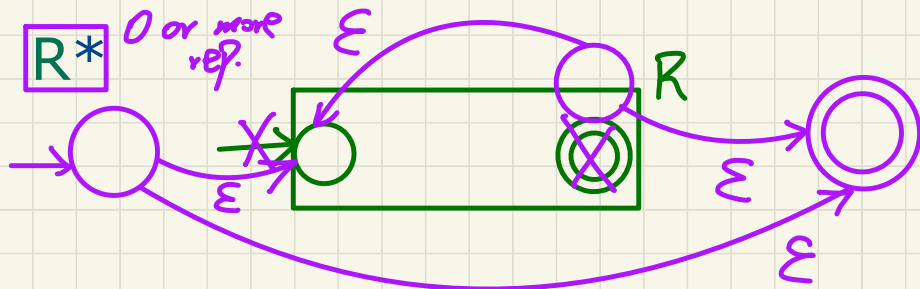


start state of left-most  $\epsilon$ -NFA

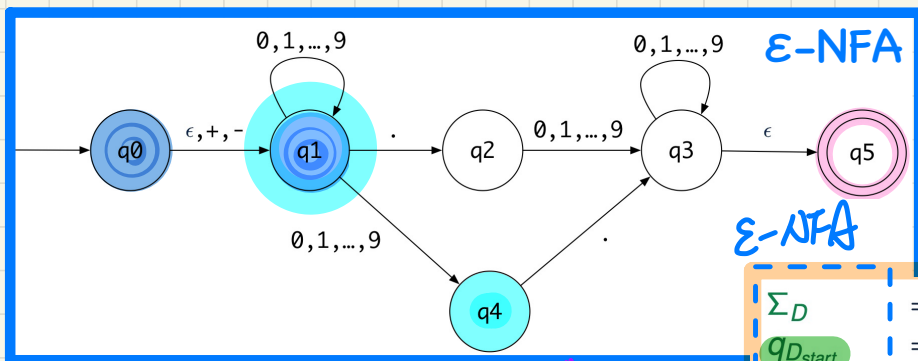
accept state of right-most  $\epsilon$ -NFA

$R^*$

0 or more rep.



# epsilon-NFA to DFA: Extended Subset Construction



$\epsilon \notin \Sigma_N$

**ε-NFA**

**DFA**

any subset state conforming in DFA is considered accepting

\*\*\*  $S \xrightarrow{a} S' \rightarrow \text{ECLOSE}(S')$  reachable states

$$\begin{aligned}
 \Sigma_D &= \Sigma_N \\
 q_{D_{start}} &= \text{ECLOSE}(q_0) \\
 F_D &= \{ S \mid S \subseteq Q_N \wedge S \cap F_N \neq \emptyset \} \\
 Q_D &= \{ S \mid S \subseteq Q_N \wedge (\exists w \bullet w \in \Sigma^* \Rightarrow S = \hat{\delta}_N(q_0, w)) \} \\
 \delta_D(S, a) &= \bigcup \{ \text{ECLOSE}(S') \mid s \in S \wedge s' \in \delta_N(s, a) \}
 \end{aligned}$$

	$d \in 0..9$	$s \in \{+, -\}$	.
$\{q_0, q_1\}$	$\{q_1, q_4\}$	$\{q_1\}$	$\{q_2\}$
$\{q_1, q_4\}$	$\{q_1, q_4\}$	$\emptyset$	$\{q_2, q_3, q_5\}$
$\{q_1\}$	$\{q_1, q_4\}$	$\emptyset$	$\{q_2\}$
$\{q_2\}$	$\{q_3, q_5\}$	$\emptyset$	$\emptyset$
$\{q_2, q_3, q_5\}$	$\{q_3, q_5\}$	$\emptyset$	$\emptyset$
$\{q_3, q_5\}$	$\{q_3, q_5\}$	$\emptyset$	$\emptyset$

**DFA**

\*  $S$  is a subset state, which can be a state in the DFA

\*\*  $S$  is reachable from  $q_0$  via  $w$  (w.r.t.  $\delta$ )

# Minimizing DFA: Algorithm

① What if  $M' = M \Rightarrow$  no optimization was necessary

② What if

$$|Q_{M'}| > |Q_M|?$$

should not be possible!

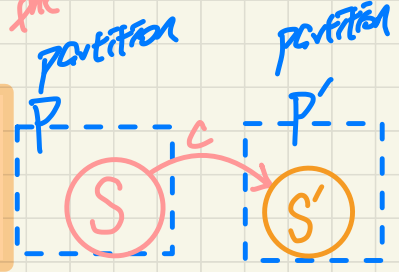
ALGORITHM: MinimizeDFAStates  
 INPUT: DFA  $M = (Q, \Sigma, \delta, q_0, F)$  *accept states*  
 OUTPUT:  $M'$  s.t. minimum  $|Q|$  and equivalent behaviour as  $M$   
 PROCEDURE:  
 $P := \emptyset$  /\* refined partition so far \*/  
 $T := \{q \in Q \mid q \in F\}$  /\* last refined partition \*/  
 while  $P \neq T$  *non-accept states*  
      $P := T$   
      $T := \emptyset$   
     for  $(p \in P)$  :  
         find the maximal  $S \subset p$  s.t. **splittable**( $p, S$ )  
         if  $S \neq \emptyset$  then  
              $T := T \cup \{S, p - S\}$   
         else  
              $T := T \cup \{p\}$   
     end

exit when reaching a fix point:  $P = T$   
 exit:  $P = T$  result from current it.  
 exit:  $P = T$  result from last it.

AS SOON AS  $P = T$  no further optimization can be done.  
 example each fix point from last it.

given a set  $P$  of states as a partition, can we find a proper subset of that can be split from the rest of  $P$ .

- splittable**( $p, S$ ) holds iff there is  $c \in \Sigma$  s.t.
- $S \subset p$  (or equivalently:  $p - S \neq \emptyset$ )
  - Transitions via  $c$  lead all  $s \in S$  to states in **same partition**  $p_1$  ( $p_1 \neq p$ ).



## Partitions of States

subsets  
(non-empty)

e.g.,  $Q = \{s_0, s_1, s_2, s_3\}$

$$P''' = \{ \{s_0\}, \{s_1, s_2\}, \{s_3\} \}$$

↳ 3 partitions.

$P'$

- Smallest number of partitions

$P''$

- Largest number of partitions

$P'''$

- Partitions somewhere in-between

- Analogy from Software Testing: Equivalent Classes

$$P' = \{Q\} = \{ \{s_0, s_1, s_2, s_3\} \}$$

↳ a single partition.

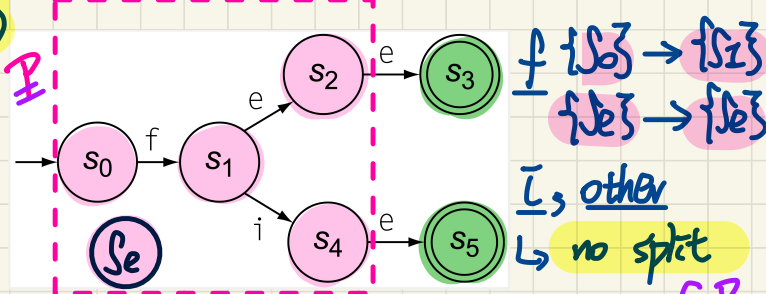
$$P'' = \{ \{s_0\}, \{s_1\}, \{s_2\}, \{s_3\} \}$$

↳ 4 partitions.

# Minimizing DFA: Example (1)

input symbols to focus on:  $f, \bar{i}, e$ , other

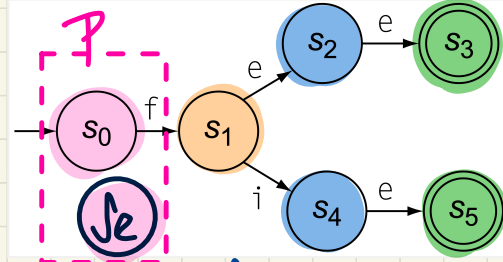
(1)



$P = \{s_0, s_1\}$   
 $\{s_0, s_1\} \xrightarrow{e} \{s_2\}$  *splittable*

$\{s_2, s_4\} \xrightarrow{e} \{s_3, s_5\}$   
 $\{s_0, s_1, s_2\} \rightarrow$

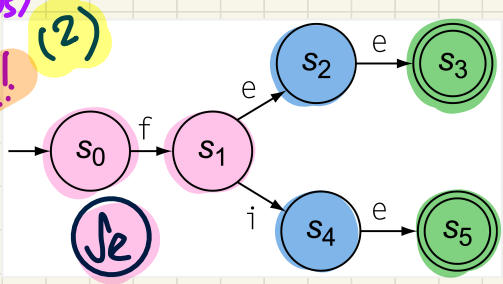
(3)



$f \{s_0\} \xrightarrow{f} \{s_1\}$  *splittable*  
 $\{s_1\} \xrightarrow{e} \{s_2\}$  *not splittable*

$\{s_2, s_4\} \xrightarrow{e} \{s_3, s_5\}$   
 $\{s_0, s_1, s_2\} \rightarrow$

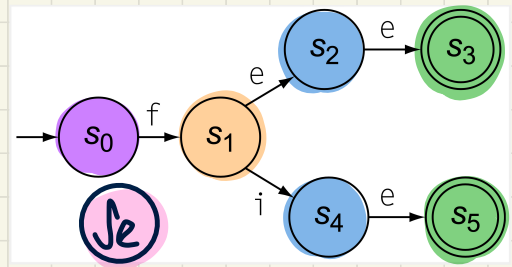
(2)



$f: \text{no split}$   
 $\bar{i}: \text{splittable}$   
 $\{s_1\} \xrightarrow{\bar{i}} \{s_4\}$

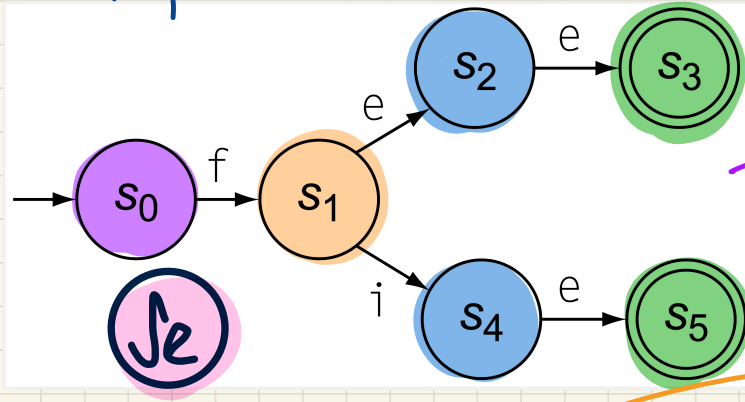
$\{s_0, s_1\} \xrightarrow{\bar{i}} \{s_4\}$   
 $\rightarrow \text{not splittable}$

(4)



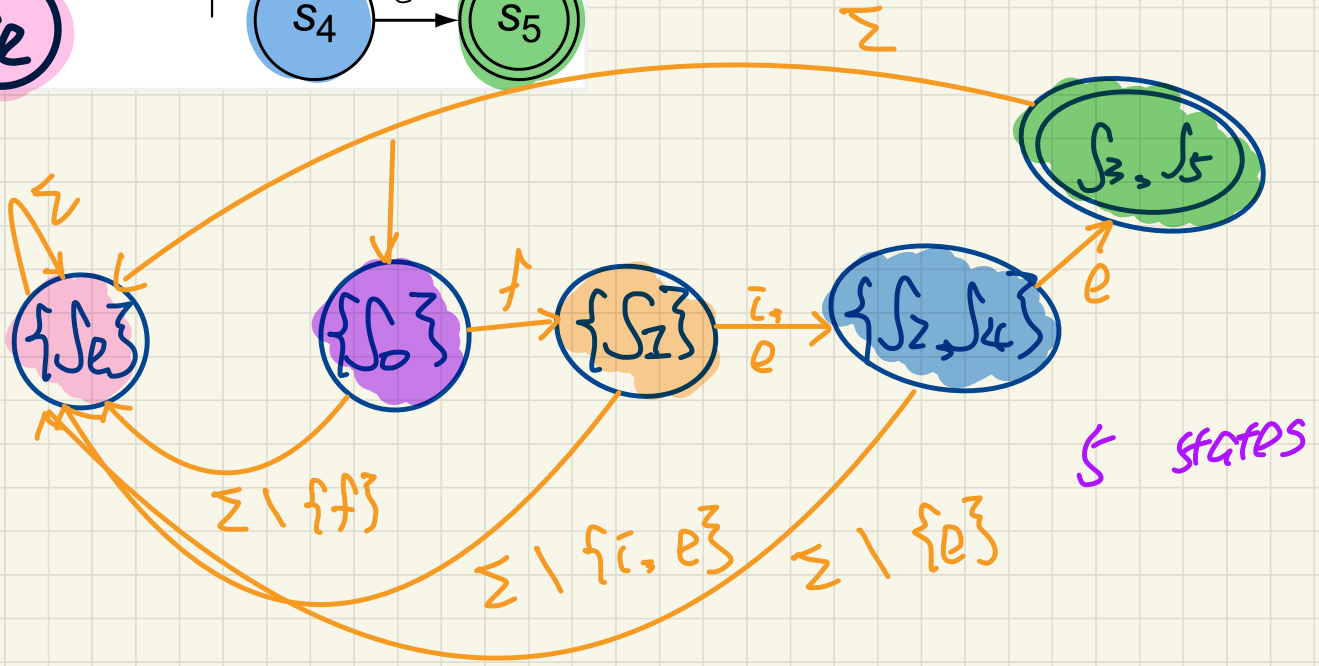
$\rightarrow \text{not splittable for any partition any move.}$

Maximal # partitions  $\rightarrow$  min # of DFA states



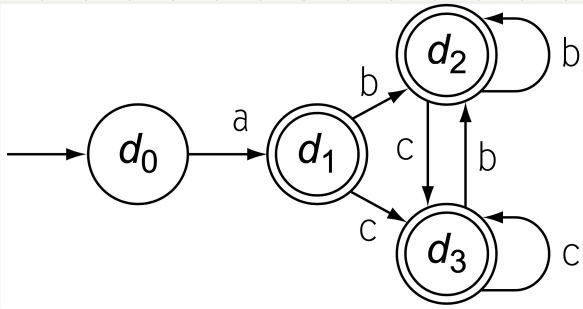
$\rightarrow$  7 states

$2/7 \approx 30\%$   
reduction.



5 states

## Minimizing DFA: Example (2)





## Minimizing DFA: Example (3)

